

Linear System Games

R. Cleve, L. Liu, W. Slofstra




AQIS 2016 August 30

Types of non-locality


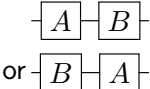
Tensor Product

Commuting Operator

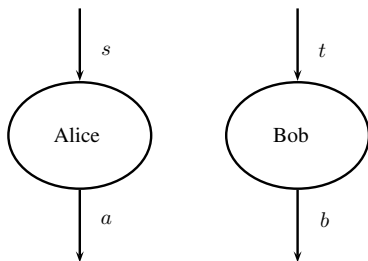
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| | Tensor Product | Commuting Operator |
|------------------|---|--------------------|
| Hilbert Space | $\mathcal{H}_A \otimes \mathcal{H}_B$ | |
| Alice's Operator | $A \otimes \mathbb{1}$ | |
| Bob's Operator | $\mathbb{1} \otimes B$ | |
| Restriction | - | |
| Circuit |  | |

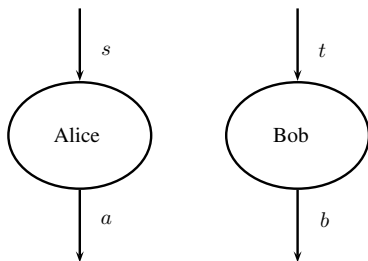
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Non-local games

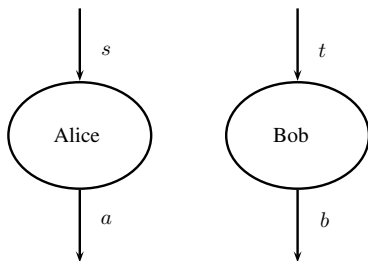


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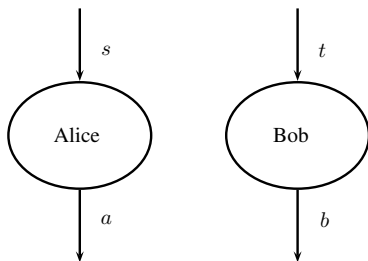
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- Output: $a, b \in \Sigma_O$

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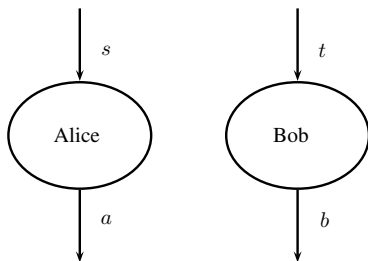
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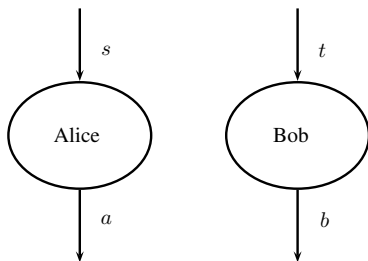


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Example: CHSH game

$a, b, s, t \in \{0, 1\}$, winning condition: $a \oplus b = st$

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Classical winning Pr: 0.75 (always output 0)

Entanglement winning Pr: $\cos^2 \frac{\pi}{8} \approx 0.85$

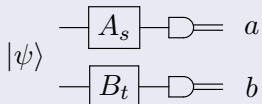
Tensor product strategies:

- Shared entangled state

$$|\psi\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B$$

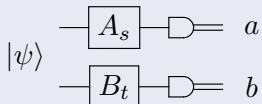
- Local operators

$$A_s \in \mathcal{H}_A \text{ and } B_t \in \mathcal{H}_B$$



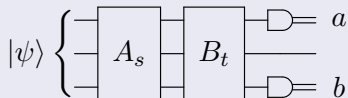
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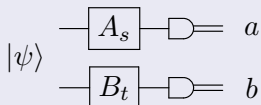
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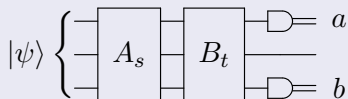
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(Classical strategy: no entanglement, no fun.)

Linear System Games

- Non-local game derived from binary¹ linear system $Ax = b$, representing set of linear constraints

¹Can be easily extended to mod p

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 - Example of a logical constraint:
 $x_1 \wedge x_2 \wedge \neg x_3 = 1$
 - Example of a linear constraint:
 $x_1 \oplus x_2 \oplus x_3 = 0$

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- Input:
 - Alice receives row $i \in \{1, \dots, m\}$, a constraint
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 - $A_{i,j} \neq 0$ (variable is in the constraint)

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- Winning Output:
 - Alice assigns vector x such that $\sum_j A_{i,j}x_j = b_i$
 - Bob outputs one bit that matches with Alice's x_j .

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Example of linear system game

- Linear system:

$$A = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \end{pmatrix}, b = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

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- Inputs: $i = 1, j = 1$
 - Alice must output x such that $x_1 \oplus x_2 = 1$
 - If Alice outputs $[1, 0, 0]$, Bob must output 1.
- Perfect classical strategy exists: setting $x = [0, 1, 0]^T$

- Classical strategy \iff classical solution

Perfect strategies

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- Tensor product strategy \iff finite dim. “operator solution”
[Cleve & Mittal 2012]

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- Commuting operator strategy \iff infinite dim. operator solution
[this paper]

Classical solution

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Alternative form

- $x_j \in \{-1, 1\}$

$$0 \mapsto 1, 1 \mapsto -1$$

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- $x_j \in \{-1, 1\}$
- $\prod_j x_j^{A_{ij}} = (-1)^{b_i}$

+ to ×

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$$x_1 x_2 = -1 \iff x_2 x_1 = -1$$

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Solution group $\langle g_1, \dots, g_n, J \rangle :$

Abstraction

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- $g_j^2 = e, J^2 = e$
- $g_j J = J g_j$

J is similar to $-\mathbb{1}$.

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\uparrow is a representation of \uparrow with $J \mapsto -\mathbb{1}$

Summary of types of solutions

| | classical sol. | operator sol. | solution group |
|-----------------------|-------------------------------------|--|--|
| binary | $x_j = \pm 1$ | $X_j^2 = \mathbb{1}$ | $g_j^2 = e, J^2 = e$ |
| constraint satisfying | $\prod_j x_j^{A_{ij}} = (-1)^{b_i}$ | $\prod_j X_j^{A_{ij}} = (-\mathbb{1})^{b_i}$ | $\prod_j g_j^{A_{ij}} = J^{b_i}$ |
| local compatibility | - | $X_j X_{j'} = X_{j'} X_j$ $A_{ij'} = A_{ij'} = 1$ | $g_j g_{j'} = g_{j'} g_j$ $A_{ij'} = A_{ij'} = 1$ |
| other | - | - | $[J, g_j] = 1 \forall g_j$ |

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- 3 The solution group satisfies $J \neq e$.

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$$1 \subseteq 2 \subseteq 3 \subseteq 4$$

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Strong Tsirelson's Conjecture (False [Slofstra 2016])

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Conne's Embedding Conjecture (Open)

- Existence of perfect solution can be ruled out by “substitution method”.

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- Not clear how to apply group structure for systems other than linear system games.

Q&A

Thank you!

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- $|\psi\rangle = (|e\rangle - |J\rangle)/\sqrt{2}$.