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# Entanglement-Assisted Quantum Communication Beating the Quantum Singleton Bound

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# Overview

- parameters of quantum error-correcting codes (QECC)
- no-cloning bound
- quantum Singleton bound
- entanglement-assisted QECC
- noiseless teleportation
- noiseless teleportation
- our scheme
- examples
- conclusions

# Code Parameters

**Quantum Code:**  $\mathcal{C} = \llbracket n, k, d \rrbracket_q$

$q$ : dimension of the subsystems

$n$ : number of subsystems used in total

$k$ : number of (logical) subsystems encoded

$d$ : “minimum distance”

– correct all errors acting on at most  $(d - 1)/2$  subsystems

– detect all errors acting on less than  $d$  subsystems

– correct all errors at less than  $d$  *known* positions (“erasures”)

**trivial absolute bound:**  $k \leq n - (d - 1) = n + 1 - d$

equivalently:  $d \leq n + 1 - k$

# No-Cloning Bound

**Assumption:**  $\mathcal{C} = \llbracket n, 1, n/2 + 1 \rrbracket_q$  exists

encoded state:

$$\sum_i \alpha_i |\psi_i\rangle |\phi_i\rangle$$

splitting:

$$\sum_i \alpha_i^2 |\psi_i\rangle \langle \psi_i|$$

$$\sum_i \alpha_i^2 |\phi_i\rangle \langle \phi_i|$$

padding:

$$\left( \sum_i \alpha_i^2 |\psi_i\rangle \langle \psi_i| \right) \otimes (|0\rangle \langle 0|)^{\otimes n/2}$$

$$(|0\rangle \langle 0|)^{\otimes n/2} \otimes \left( \sum_i \alpha_i^2 |\phi_i\rangle \langle \phi_i| \right)$$

correction:

$$\sum_i \alpha_i |\psi_i\rangle |\phi_i\rangle$$

$$\sum_i \alpha_i |\psi_i\rangle |\phi_i\rangle$$

two independent copies

$\implies$  no-cloning bound:  $d - 1 < n/2$  or  $2d < n + 2$

# Quantum Singleton Bound

[E. Rains, Nonbinary Quantum Codes, quant-ph/9703048]

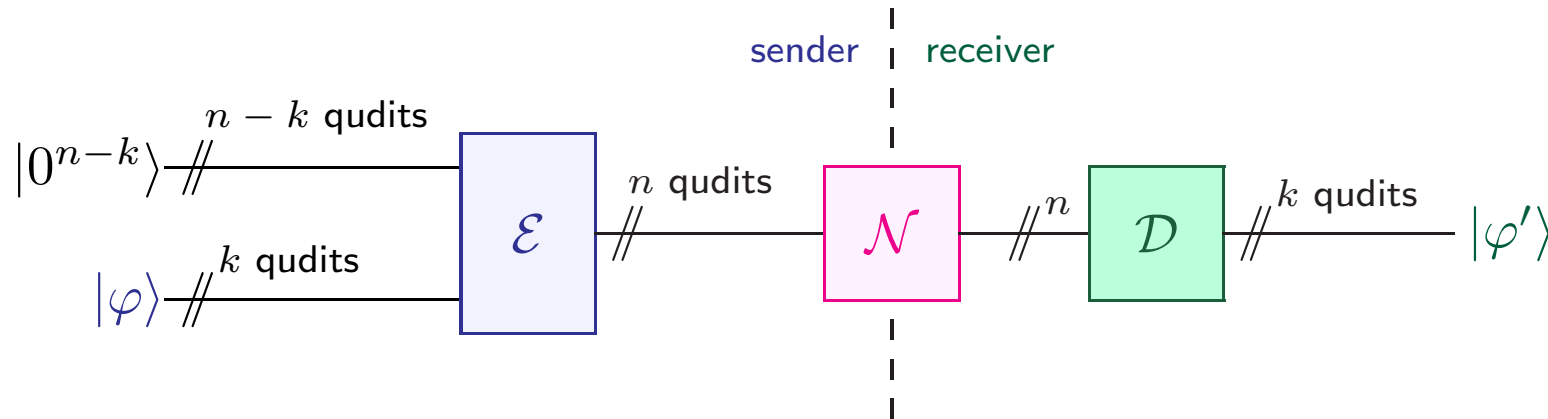
Let  $\mathcal{C} = \llbracket n, k, d \rrbracket_q$  be a quantum error-correcting code. Then

$$2d \leq n + 2 - k.$$

**Proof (sketch:)** [J. Preskill, Lecture Notes, Sect. 7.8.3]

- splitting of the code into  $(d - 1) + (d - 1) + (n + 2 - 2d)$  qudits
- maximally entangled state of input and reference system  $R$
- the first two blocks with  $d - 1$  qudits have no correlations with  $R$
- the entropy of the third block with  $n + 2 - 2d$  qudits is at least  $k \log q$

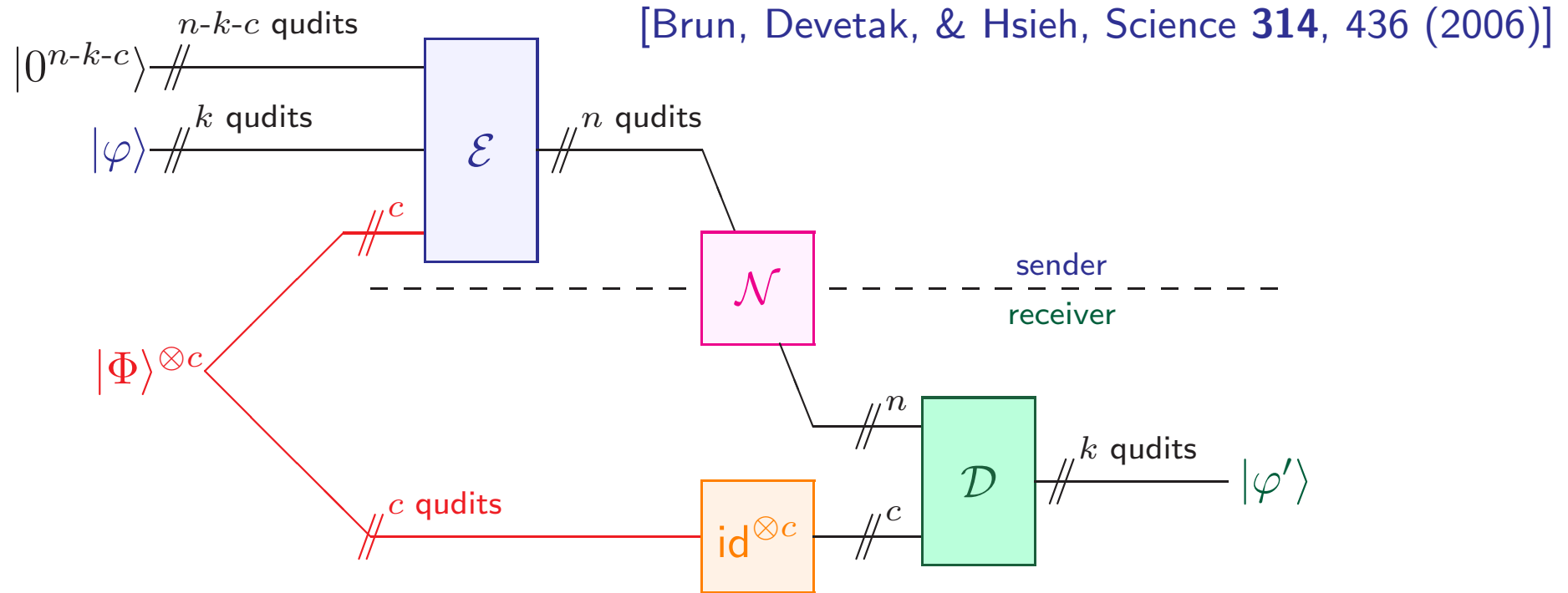
# Quantum Error-Correcting Code (QECC)



Scheme of a communication protocol using a QECC  $[[n, k, d]]_q$

quantum Singleton bound:  $2d \leq n + 2 - k$

# Entanglement-Assisted QECC



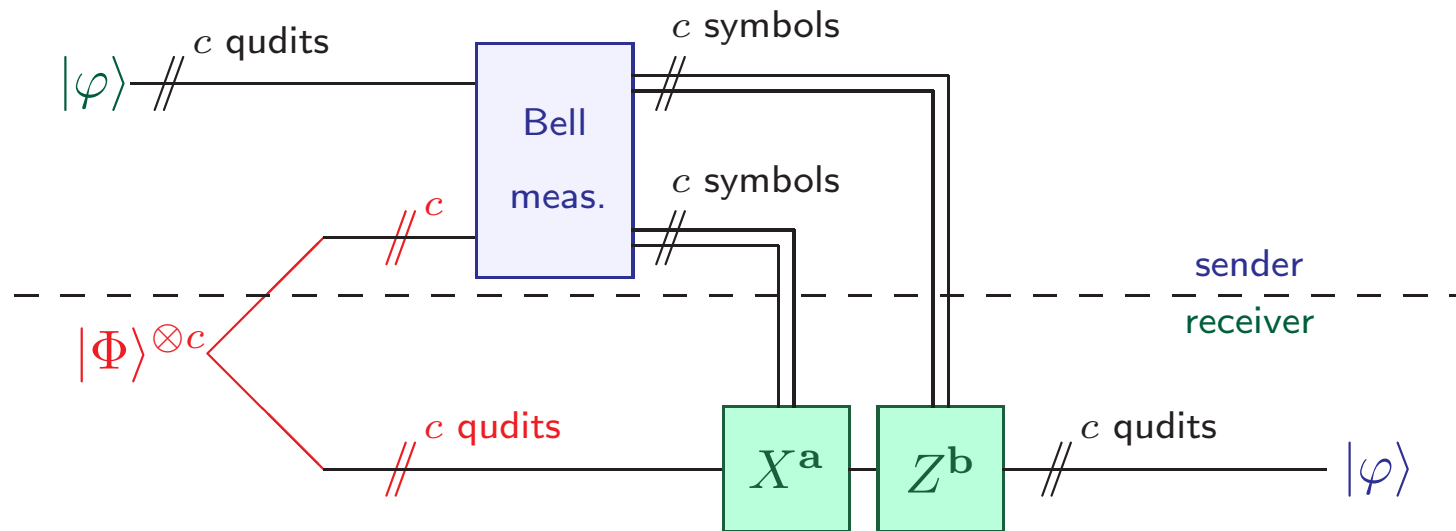
Communication scheme using an entanglement-assisted QECC  $[[n, k, d; c]]_q$ .

*similar* to a QECC of length  $n + c$

$\implies$  **conjectured bound:**  $2d \leq n + c + 2 - k$

[Brun, Devetak, & Hsieh, arXiv:quant-ph/0608027]

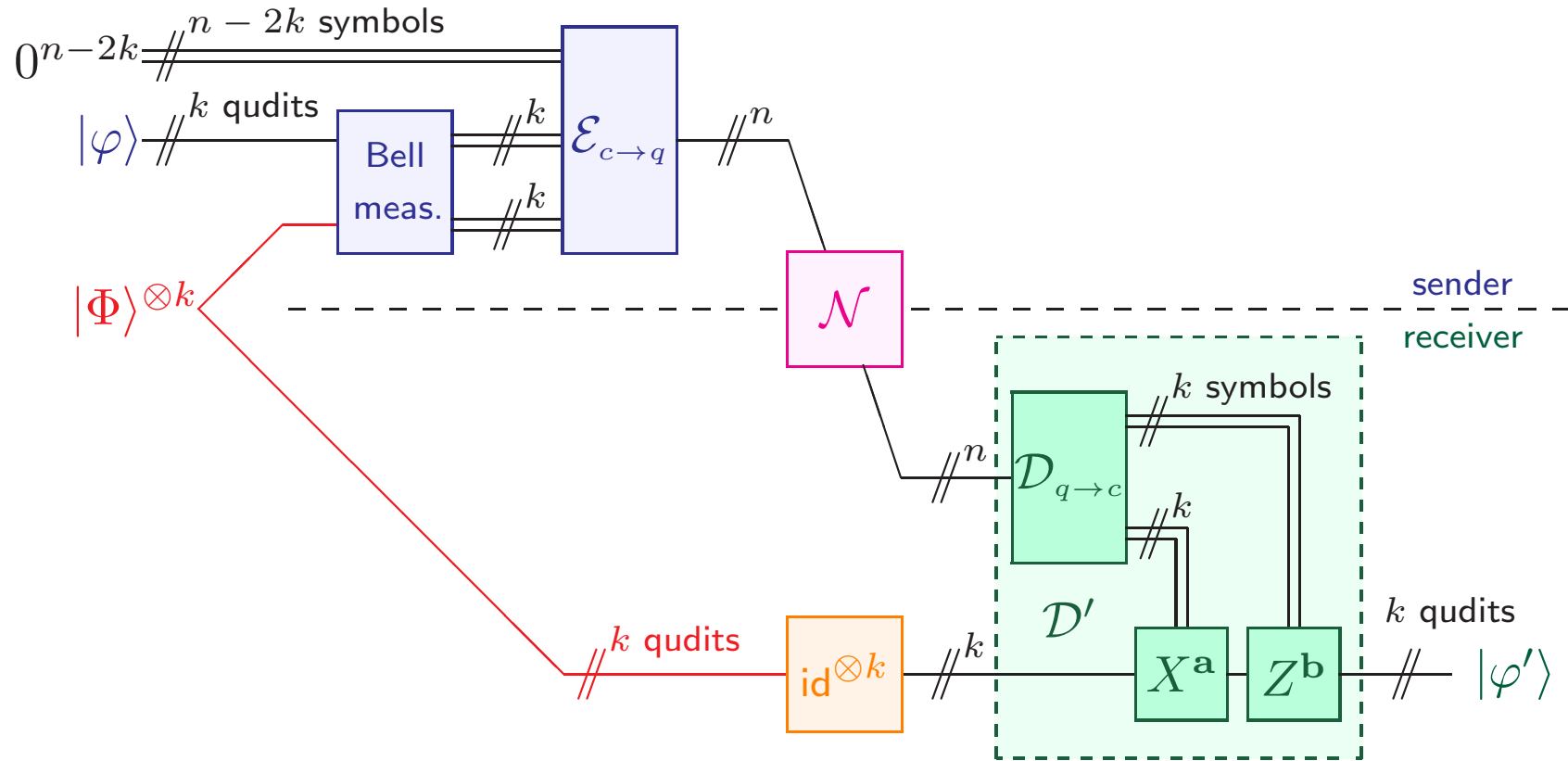
# Noiseless Teleportation



The maximally entangled states  $|\Phi\rangle^{\otimes c}$  can be prepared by either party, or even a third party.



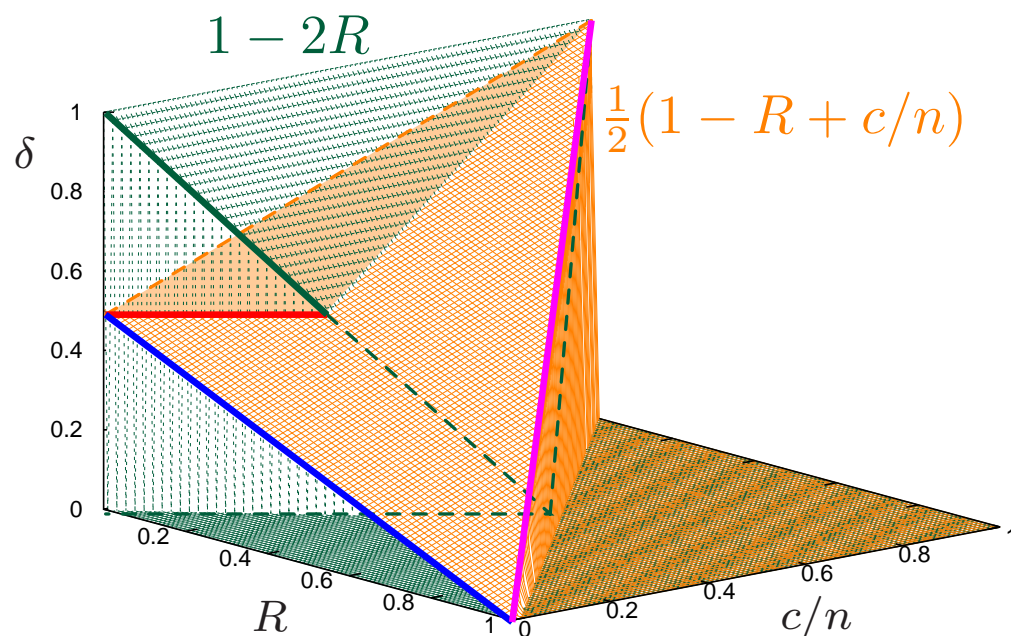
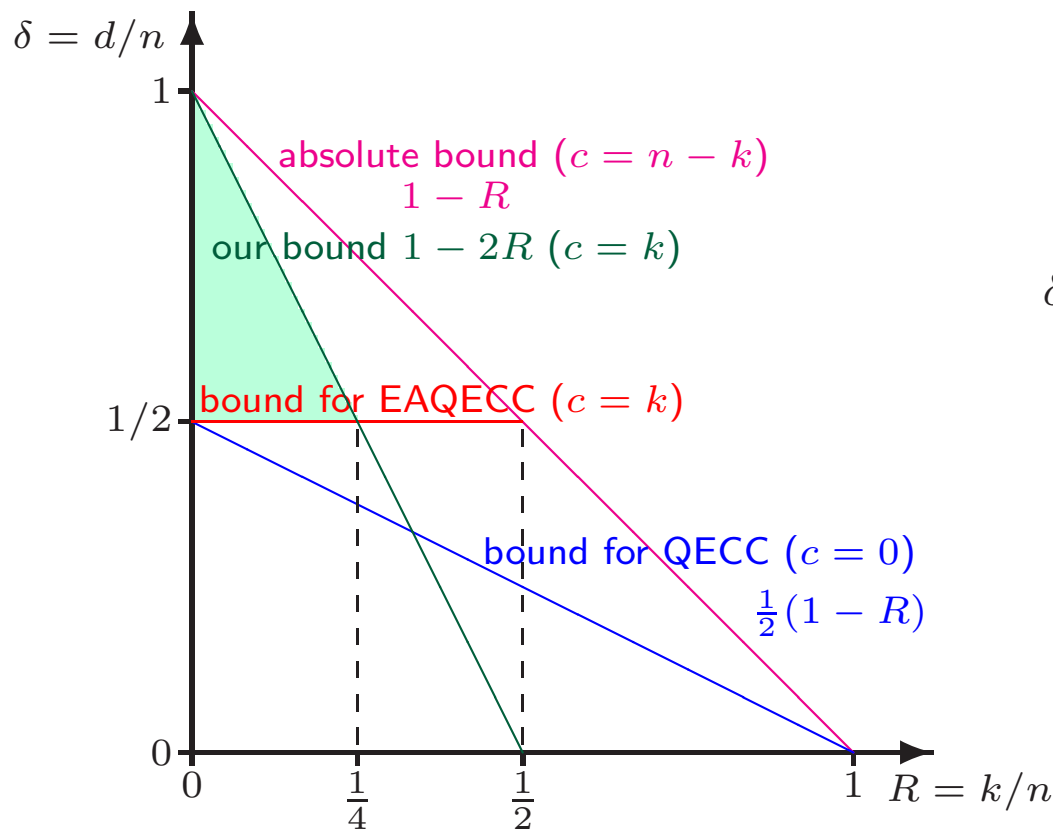
# The New Scheme



Teleportation-based scheme using  $c = k$  maximally entangled states.

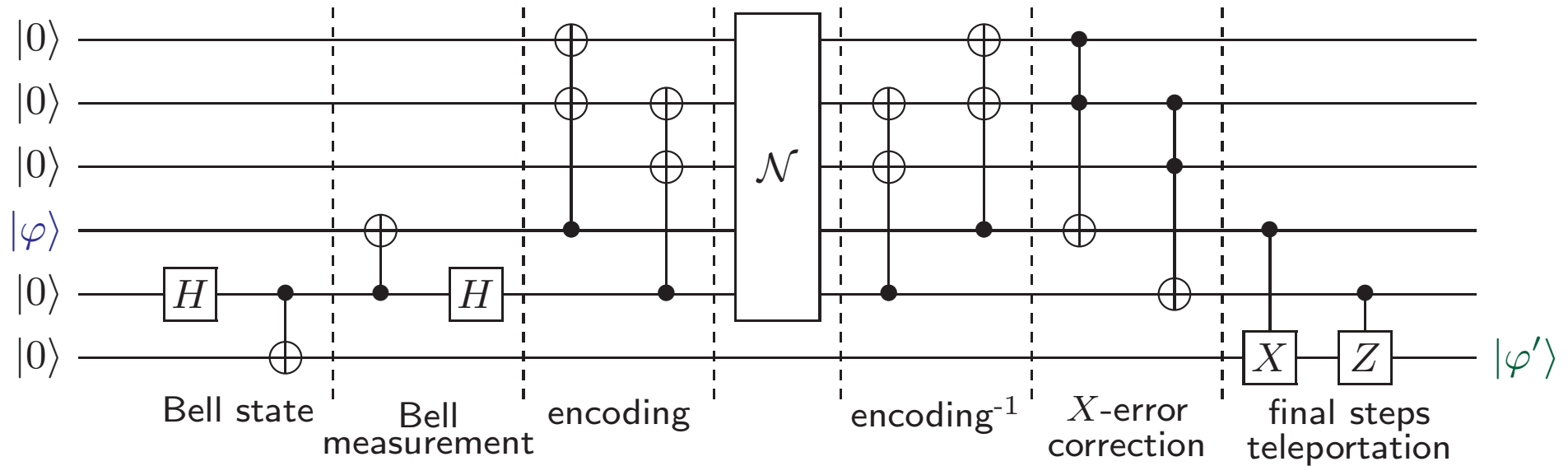
**bound:**  $d \leq n + 1 - 2k$

# Bounds



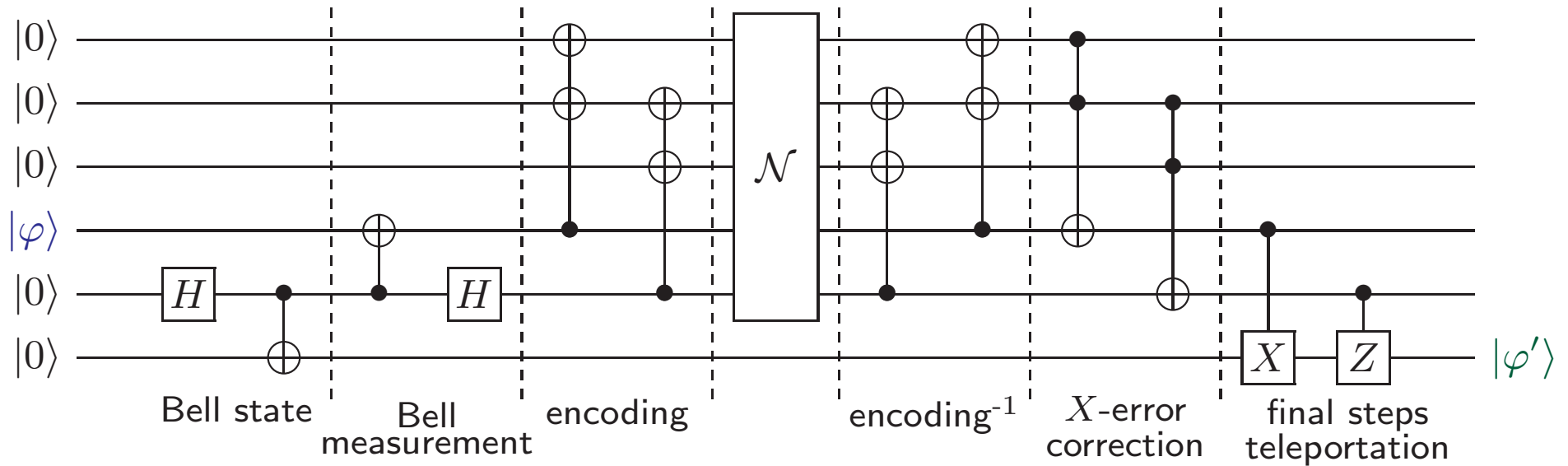
Asymptotic bounds (length  $n \rightarrow \infty$ ) on the normalized minimum distance  $\delta = d/n$  as a function of the code rate  $R = k/n$ .

# Example: $[[5, 1, 3; 1]]_2$



- classical code  $[5, 2, 3]_2$  generated by  $G = \begin{pmatrix} 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 \end{pmatrix}$
- standard QECC  $[[5, 1, 3]]_2$  exists, but our scheme additionally tolerates all phase errors

# Example: $[[5, 1, 3; 1]]_2$



stabilizer & normalizer:

$$\left( \begin{array}{cccccc|c}
 Z & I & I & I & I & I & I \\
 I & Z & I & I & I & I & I \\
 I & I & Z & I & I & I & I \\
 I & I & I & I & Z & I & I \\
 I & I & I & I & I & Z & I \\
 \hline
 I & I & I & X & I & I & I \\
 I & I & I & Z & I & I & I
 \end{array} \right) \rightarrow \left( \begin{array}{cccccc|c}
 Z & I & I & Z & I & I & I \\
 I & Z & I & Z & Z & I & I \\
 I & I & Z & I & Z & I & I \\
 X & X & I & X & Z & X & X \\
 I & X & X & I & X & Z & Z \\
 \hline
 X & X & I & X & I & I & I \\
 I & X & X & Z & X & I & I
 \end{array} \right)$$

## Example: $[[5, 1, 3; 1]]_2$

stabilizer & normalizer:

$$\left( \begin{array}{ccccc|c} Z & I & I & I & I & I \\ I & Z & I & I & I & I \\ I & I & Z & I & I & I \\ I & I & I & I & Z & I \\ I & I & I & I & I & Z \\ \hline I & I & I & X & I & I \\ I & I & I & Z & I & I \end{array} \right) \rightarrow \left( \begin{array}{ccccc|c} Z & I & I & Z & I & I \\ I & Z & I & Z & Z & I \\ I & I & Z & I & Z & I \\ X & X & I & X & Z & X \\ I & X & X & I & X & Z \\ \hline X & X & I & X & I & I \\ I & X & X & Z & X & I \end{array} \right)$$

split weight enumerators [C.-Y. Lai & A. Ashikhmin, arXiv:1602.00413v2]

$$W_{\mathcal{C}}^{(5,1)}(Y; V) = (1 + 2Y^2 + 4Y^3 + Y^4) + (4Y^3 + 12Y^4 + 8Y^5)V$$

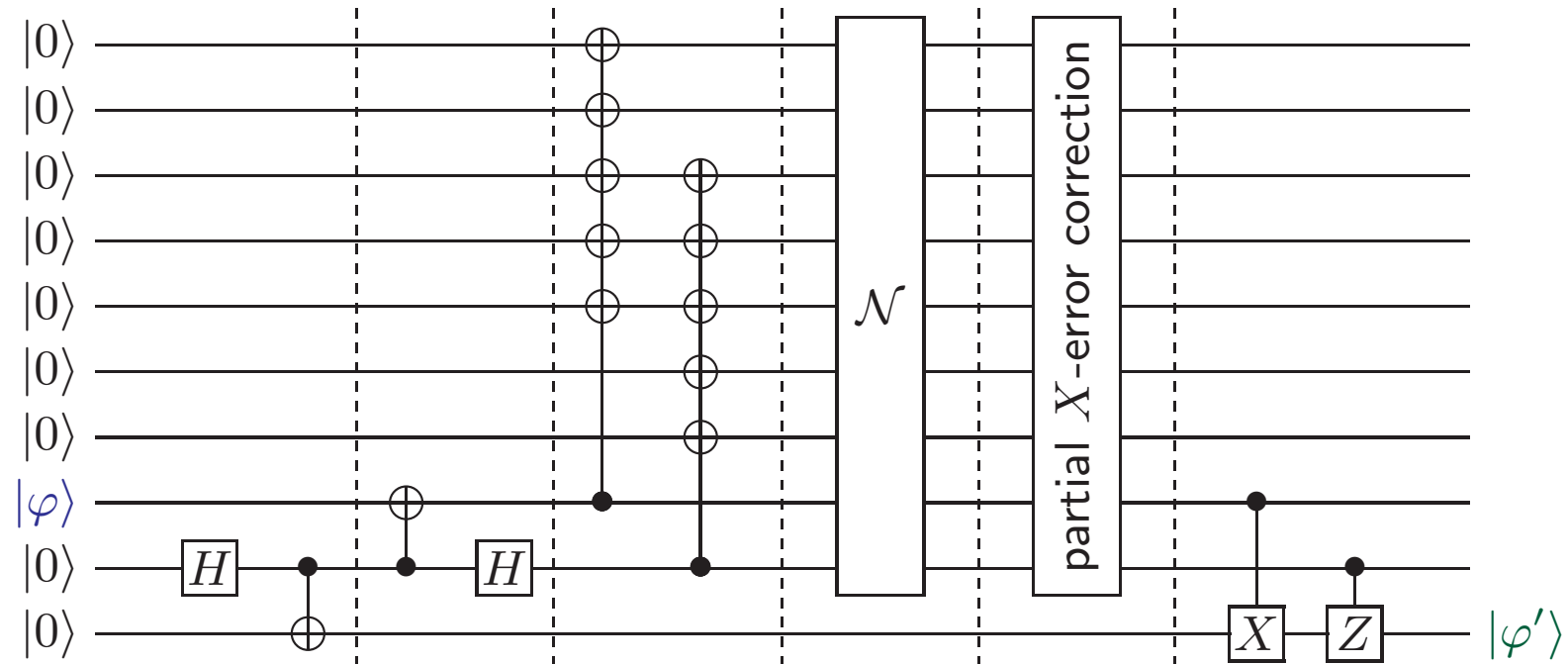
$$W_{\mathcal{C}^\perp}^{(5,1)}(Y; V) = (1 + 2Y^2 + 8Y^3 + 13Y^4 + 8Y^5)$$

$$+ (5Y + 8Y^2 + 18Y^3 + 40Y^4 + 25Y^5)V$$

$$W_{\mathcal{C}^\perp}^{(5,1)}(Y; V) - W_{\mathcal{C}}^{(5,1)}(Y; V) = (4Y^3 + 12Y^4 + 8Y^5)$$

$$+ (5Y + 8Y^2 + 14Y^3 + 28Y^4 + 17Y^5)V$$

# Example: $[[9, 1, 6; 1]]_2$



- classical code  $[9, 2, 6]_2$

- original bound for  $n = 9, k = 1, c = 1$ :  $2d \leq n + c + 2 - k = 11$

# Example: $[[9, 1, 6; 1]]_2$

stabilizer & normalizer:

$$\left( \begin{array}{cccccccc|c} Z & I & I & I & I & I & I & Z & I & I \\ I & Z & I & I & I & I & I & Z & I & I \\ I & I & Z & I & I & I & I & Z & Z & I \\ I & I & I & Z & I & I & I & Z & Z & I \\ I & I & I & I & Z & I & I & Z & Z & I \\ I & I & I & I & I & Z & I & I & Z & I \\ I & I & I & I & I & I & Z & I & Z & I \\ X & X & X & X & X & I & I & X & Z & X \\ I & I & X & X & X & X & X & I & X & Z \\ \hline X & X & X & X & X & I & I & X & I & I \\ I & I & X & X & X & X & X & Z & X & I \end{array} \right)$$

## Example: $[[9, 1, 6; 1]]_2$

split weight enumerators [C.-Y. Lai & A. Ashikhmin, arXiv:1602.00413v2]

$$W_{\mathcal{C}}^{(9,1)}(Y; V) = (1 + 9Y^2 + 27Y^3 + 27Y^4 + 27Y^5 + 27Y^6 + 9Y^7 + Y^9) \\ + (48Y^6 + 144Y^7 + 144Y^8 + 48Y^9)V$$

$$W_{\mathcal{C}_{\perp}}^{(9,1)}(Y; V) = (1 + 9Y^2 + 27Y^3 + 27Y^4 + 27Y^5 + 75Y^6 + 153Y^7 + 144Y^8 + 49Y^9) \\ + (9Y + 27Y^2 + 57Y^3 + 99Y^4 + 99Y^5 + 201Y^6 + 459Y^7 + 441Y^8 + 144Y^9)V$$

$$W_{\mathcal{C}_{\perp}}^{(9,1)}(Y; V) - W_{\mathcal{C}}^{(9,1)}(Y; V) \\ = (48Y^6 + 144Y^7 + 144Y^8 + 48Y^9) \\ + (9Y + 27Y^2 + 57Y^3 + 99Y^4 + 99Y^5 + 153Y^6 + 315Y^7 + 297Y^8 + 96Y^9)V$$



# Conclusions

## Summary

- using a simple teleportation-based scheme beats the quantum Singleton bound
- the scheme uses the same assumptions & operations as “standard” entanglement-assisted quantum codes

## Outlook

- new schemes to get more flexibility how much entanglement is used
- derive new upper bounds, e.g., using linear programming techniques