

# Characterizing Quantum Supremacy in Near-Term Devices

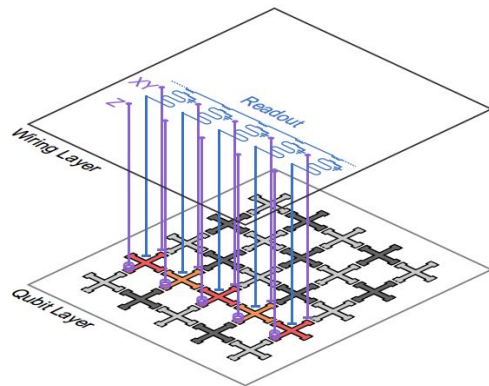
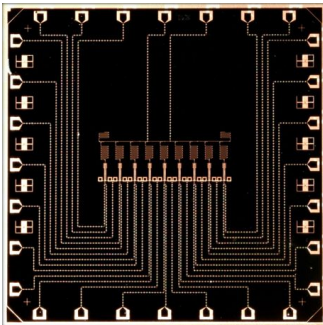
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**Google mission:** build a quantum computer that does something useful, ASAP



Aiming for medium scale planar platform (e.g. 7 x 7 qubits) in near future

Long road to error correction; can intermediate devices perform useful computation?

We will solve *useless* and intractable problems before *useful* and intractable problems

**1608.00263: minimal resource experimental demonstration of “Quantum Supremacy”**

# What is “Quantum Supremacy”?

**Preskill (1203.5813):** “when well controlled quantum systems can perform tasks surpassing what can be done in the classical world”

**Boixo et al. (1608.00263):** “when existing quantum devices can solve formal computational problems that cannot be solved on existing classical devices by implementing known algorithms in a reasonable amount of time”

**Relative concept:** “supremacy” implies competition with classical state-of-the-art

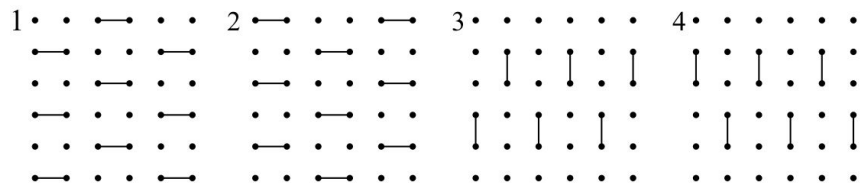
**Practical:** a description of status quo, not complexity theory result, not asymptotic

# Task: sample from the output of a random quantum circuit

**Select RQC  $U$ , apply to  $|+\rangle^n$  state, sample bit strings  $|x\rangle$  in computational basis**

Produce samples with high likelihood,  $\prod_m p_U(x_m)$  where  $p_U(x) = |\langle x|U|+\rangle|^2$

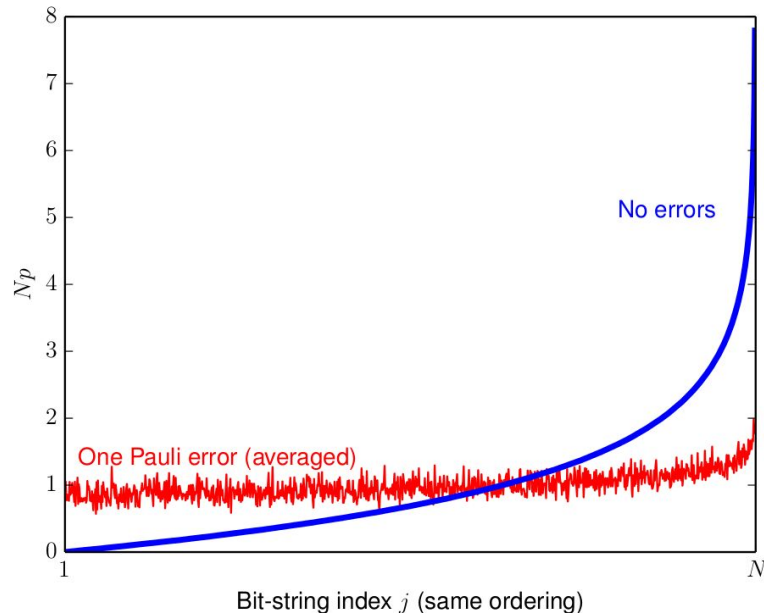
**Our RQCs consist of CZ cycles on planar lattice with random  $X^{1/2}$ ,  $Y^{1/2}$  and T gates**



**RQCs converge to finite moments of Haar measure**

Even single qubit Pauli errors decorrelate samples

**Classical cost exponential in  $\min(n, d n^{1/2}, \tau)$**



# A practical method to estimate the fidelity of RQCs

**Cross entropy  $H_{U,A} \equiv -\sum_x p_A(x) \log p_U(x)$  measures similarity of  $p_U(x)$  and  $p_A(x)$**

**$p_U(x) = |\langle x|U|+\rangle|^2$  follows Porter-Thomas distribution:  $\text{Prob}[p_U(x) > p] = N e^{-Np}$**

When  $p_A(x) = p_U(x)$  then  $H_{U,A} = -\int p \log p (N e^{-Np}) N dp = \log N - 1 + \gamma \equiv H_1$

For average  $U$ , when  $p_A(x)$  uncorrelated with  $p_U(x)$  then  $\langle H_{U,A} \rangle_U = H_1 + 1 \equiv H_0$

**Cross entropy difference is  $\alpha_A \equiv H_0 - \langle H_{U,A} \rangle_U$  and approximates fidelity for RQCs**

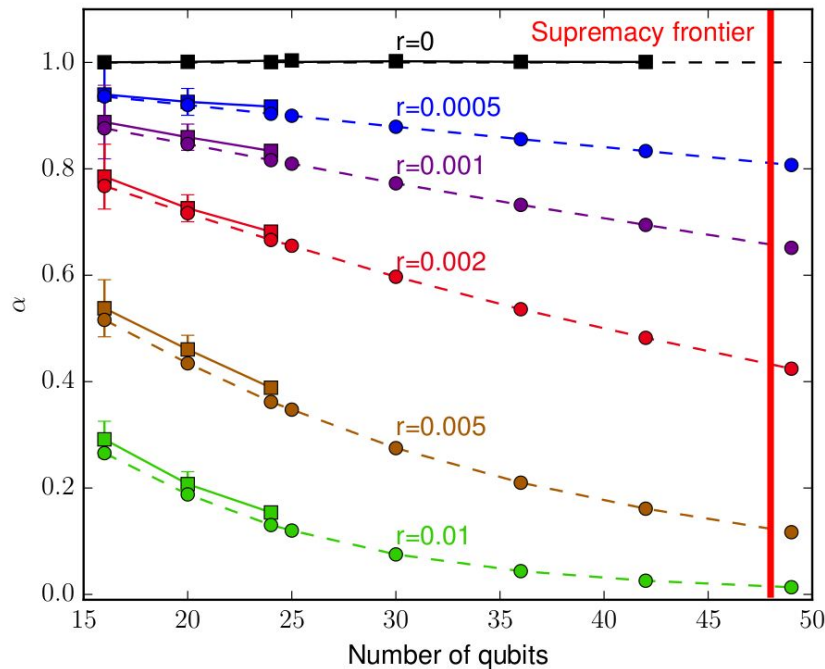
Consider noisy quantum circuit with fidelity  $f$  so that  $p_A(x) = f p_U(x) + (1 - f) p_U^\varepsilon(x)$

$\alpha_A = H_0 + f \sum_x p_U(x) \log p_U(x) + (1 - f) \langle \sum_x p_U^\varepsilon(x) \log p_U(x) \rangle_U \approx H_0 - f H_1 - (1 - f) H_0 = f$

**If  $p_U(x)$  computable, CED efficiently measured as  $H_0 + m^{-1} \sum_k \log p_U(x_k^A) + O(m^{-1/2})$**

# The experimental proposal for quantum supremacy

1. Randomly select RQCs, experimentally implement  $U|+\rangle^n$ ; sample  $m$  strings  $\{x_1, x_2, \dots, x_m\}$  in computational basis
2. Use classical supercomputer to compute probabilities  $\{p_U(x_1), p_U(x_2), \dots, p_U(x_m)\}$  then,  
$$f = \alpha = \log N + \gamma + m^{-1} \sum_k \log p_U(x_k) + O(m^{-1/2})$$
3. Extrapolate  $\alpha$  to supremacy regime;  
expect  $f \propto \exp(-r \cdot \text{number of gates})$



**Probabilistic verification procedure possible**

# Summary of results

**Planar RQCs with  $d > 27$  and  $n = 7 \times 7$  likely cost  $\Omega(2^n)$  to classically simulate**

Would require at least 4 Petabytes of RAM with fast interconnect (doesn't exist)

**Cross entropy difference is useful metric for benchmarking quantum circuits**

Approximates fidelity for RQCs, efficient to measure if classical simulation possible

**Sampling RQCs and achieving target CED defines minimal supremacy experiment**

Based on current fidelities and digital error model, supremacy is within reach

**Provided complexity theoretic reasons to believe this task is fundamentally hard**

Extended results from IQP circuits (similar to BosonSampling) to broader class

# Acknowledgements

## Coauthors

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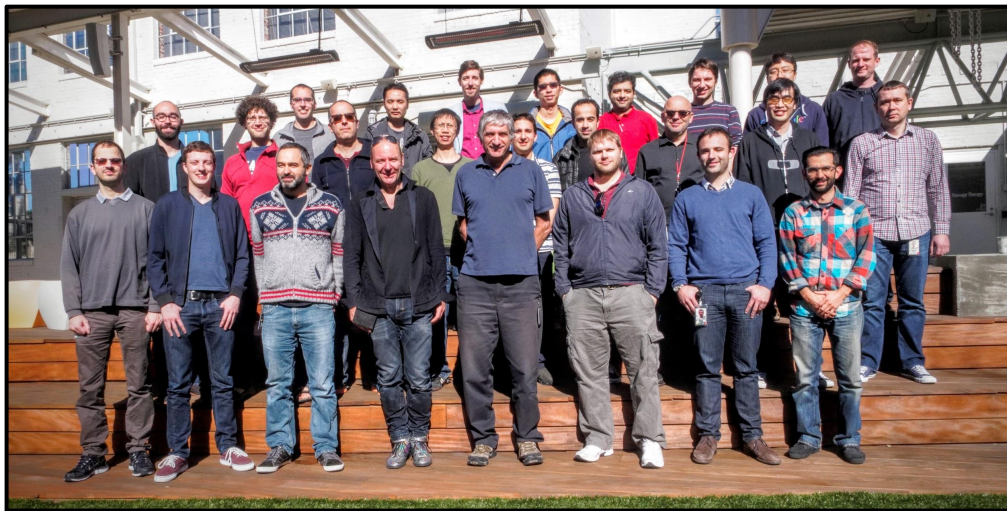
John Martinis (Google)

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## Large circuit simulations by

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Jeff Hammond (Intel)



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# Complexity theoretic implications

$p_U(x) = |\langle x|U|+\rangle|^2 = \lambda |Z|^2$  **computable using path-integral representation of U**

For IQP circuits, Z is Ising partition function  $Z = \sum_s e^{-i H(s)}$  with  $H(s) = h s + s J s$

**Z has sign problem; conjectured that Z is not approximable with NP-oracle**

If  $p_U(x)$  efficiently sampled, Z is approximable with NP-oracle (1504.07999)

**Implies that efficient classically sampling would collapse polynomial hierarchy**

Old result for BosonSampling (1011.3245) and IQP circuits (1005.1407)

**We follow methodology of 1504.07999 and show same implications**

Whereas IQP circuits map to 2D Ising model, our circuits map to 3D Ising model

