A semi-device-independent framework based on natural physical assumptions and its application to random number generation

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There exist vulnerabilities in quantum cryptography, successfully exploited by quantum hackers

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We explore bright-light control of superconducting nanowire single-photon detectors (SNSPDs) in the shunted configuration (a practical measure to avoid latching). In an

These attacks exploit a mismatch between the theoretical model used to prove security and the actual implementation

Device-independent quantum cryptography



Region not allowed by quantum theoryNG and

This approach can be used to certify the security of RNG and QKD protocols, or even the performance of quantum computers.

Usual, "device-dependent" quantum cryptography

Based on a detailed characterization of the devices Semi-device-independent quantum cryptography

Based on a few assumpions. Devices are partly untrusted.

E.g.:

- Measurement-device-independence
- One-sided quantum cryptography
- Source-independent QRNG
- Qubit assumption
- Source & measurement independence
- ...

Advantage: higher rate, easier to implement than fully device-independent protocols

Fully "device-independent" quantum cryptography

Based on minimal assumptions. Devices can be untrusted.



Hopefully, it can also be used for QKD

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- Energy constraint

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Outline

- Why semi-device-independent quantum cryptography?
- Motivation for our energy constraint assumption
- Results

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Even full DI requires non-trivial assumptions



DI RNG implementations

Monroe experiment





we extracted 256 bits, certified to be uniform to within 0.001.

in [18], which is titled "XOR 3" and consists of a total of 182, 161, 215 trials, acquired in 30 min of running the experiment, improving on the approximately one month duration of data





Don't waste time developing cars: in the future planes will be easy to build, common, and affordable.



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Semi-DI protocols based on a qubit assumption



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Mismatch between model used for security proof and implementation!



Qubit assumption is an idealization.
 → Shows that it is important to choose well the assumptions.







Click with prob $1 - \alpha^2$ No-click with prob α^2



$\alpha|0\rangle+\beta|1\rangle+\gamma|2\rangle+\cdots$





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Preparation

Measurement



Assumption: $\langle 0 | \rho_x | 0 \rangle \ge \omega_x$





- Natural relaxation of the no-communication assumption of full DI protocols



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- The appropriate space to describe quantum optics experiments is the Fock space of several quantum optical modes. In this context, it is natural to bound the average number of photons.
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- This is an assumption anyway made in many quantum optics experiments in which attenuated laser sources are used.
- It is directly related to simple characteristics of the device components (laser power, attenuator) and robust to device imperfections.
- It could be directly monitored (calibrated power meter) or enforced (optical fuse).

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No-communication assumption



Qubit assumption



Violation of Bell inequalities Q > C Violation of "dimension witnesses" Q > C

Energy constraint assumption



x = 1, 2



Input-output statistics

 $P(b|x) = Tr[\rho_x M_b]$ or $P(b|x) = \sum_{\lambda} p_{\lambda} Tr[\rho_x^{\lambda} M_b^{\lambda}]$

Equivalent to knowledge of the bias of *b* given *x*: $E_x = P(b = 1|x) - P(b = -1|x)$ $E_x = Tr[\rho_x M] \text{ or } E_x = \sum_{\lambda} p_{\lambda} Tr[\rho_x^{\lambda} M^{\lambda}]$

Output of devices is non-trivial if b is correlated to x Amount of correlations can be measured by quantity $E_{-} = (E_{1}-E_{2})/2$ Probability to guess correctly x given b is $\frac{1}{2} + \frac{|E_{-}|}{2}$

- *b* does not depend on $x: E_{-} = 0$
- *b* fully correlated to $x: |E_-| = 1$

x = 1, 2



Assumption

Energy constraint: $\langle 0 | \sum_{\lambda} p_{\lambda} \rho_{x}^{\lambda} | 0 \rangle \ge w_{x}$

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$$\sum_{\lambda} p_{\lambda} Tr[H\rho_{x}^{\lambda}] \leq 1 - w_{x}$$

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Maximal value for E_{-} given $w_{1} = w_{2} = w$?



•
$$w = 1 \rightarrow E_{-} = (E_1 - E_2)/2 = 0$$

•
$$w = 1/2 \rightarrow |\rho_{1,2}\rangle = (|0\rangle \pm |1\rangle)/\sqrt{2}$$

 $\rightarrow E_{-} = (E_{1} - E_{2})/2 = 1$

•
$$\frac{1}{2} \le w \le 1$$
 arbitrary
 $|\rho_1\rangle = \sqrt{w}|0\rangle + \sqrt{1-w}|\phi_1\rangle$
 $|\rho_2\rangle = \sqrt{w}|0\rangle + \sqrt{1-w}|\phi_2\rangle$

Scalar product minimal if $|\phi_1\rangle = -|\phi_2\rangle = |1\rangle$ $\Rightarrow |\rho_{1,2}\rangle = \sqrt{w}|0\rangle \pm \sqrt{1-w}|1\rangle$

Best distinguishing measurement: $M = \sigma_{\chi}$

→ We find the inequality
$$E_{-} = \frac{E_1 - E_2}{2} \le 2\sqrt{w(1 - w)}$$

According to quantum strategies: $E_{-} = \frac{E_1 - E_2}{2} \le 2\sqrt{w(1 - w)} = Q_{\text{max}}$ Maximal value for "classical" strategies?

How to define "classical" strategies? One possibility:

> "classical" strategies = "deterministic" strategies (or convex combinations thereof) $E_x = \sum_{\lambda} p_{\lambda} E_x^{\lambda}$ with $E_x^{\lambda} = \pm 1$

Let's be more conservative and compare Q_{\max} to strategies where only E_1 is deterministic $E_1 = \sum_{\lambda} p_{\lambda} E_1^{\lambda}$ with $E_1^{\lambda} = \pm 1$, no constraint on E_2

→ If $Q_{\max} > D_{\max}$ → the output of x = 1 is random (even to adversary with arbitrary knowledge of the devices)

$$E_1 = \langle \rho_1 | M | \rho_1 \rangle = 1 \implies M = 2 | \rho_1 \rangle \langle \rho_1 | - I$$
$$\implies E_2 = \langle \rho_2 | M | \rho_2 \rangle = 2 | \langle \rho_2 | \rho_1 \rangle |^2 - 1$$
$$\implies E_- = \frac{E_1 - E_2}{2} = 2 - 2 | \langle \rho_2 | \rho_1 \rangle |^2$$

Minimal value of $|\langle \rho_2 | \rho_1 \rangle|^2$ given w

$$\Rightarrow E_{-} \leq 4w(1-w) = D_{\max}$$





Energy constraint: $\langle 0 | \sum_{\lambda} p_{\lambda} \rho_{x}^{\lambda} | 0 \rangle \ge w_{x}$

•
$$E_{-} = \frac{E_1 - E_2}{2}$$

- If E_1 is deterministic, we have the "Bell inequality" $E_- \le 4w(1-w) = D_{\max}$
- According to a general quantum strategy $E_{-} \leq 2\sqrt{w(1-w)} = Q_{\max} = \sqrt{D_{\max}} > D_{\max}$

More generally, it is possible to characterize completely the set of allowed values (E_1, E_2) for given energy bounds (w_1, w_2)



One nice way to do it:

Given this characterization, one can also put rigorous bounds on the output entropy given (E_1, E_2) \rightarrow straightforward to build a RNG protocol where amount of randomness produced is evaluated assuming only the energy bound, but no other assumption on the devices. How to produce "non-deterministic" correlations in the lab?



A practical implementation with gaussian states and homodyne measurements:



Homodyne measurement of X quadrature with b = sign(X)

A simpler implementation with a slightly stronger assumption



Summary

- We propose to use a bound on the energy of optical signals as a unique assumption on which to prove the security of prepare-and-measure quantum cryptography protocol (with no other assumptions on the devices)
- We have shown that there is a gap between what can be achieved with very simple quantum implementations and deterministic strategies. This is equivalent to the violation of Bell inequalities in full DI protocols.
- These results immediately imply the existence of RNG protocols where the amount of randomness produced can be certified without making any assumptions about the devices except the energy assumption.

Open question

- Is the energy assumption sufficient to prove the security of a QKD protocol?
- We implicitly assumed that the preparation and measurement device did not share prior entanglement. Can this be relaxed?
- One extra motivation for the energy assumption is that it is in principle compatible with CV protocols for which no DI or semi-DI implementations have been introduced.

Can we analyze the security of a genuinely CV protocol in a DI setting using the energy assumption?